

## 编译原理 课后题答案

## 第二章

## P36-6

(1)

 $L(G_1)$  是 0~9 组成的数字串

(2)

最左推导:

 $N \Rightarrow ND \Rightarrow NDD \Rightarrow NDDD \Rightarrow DDDD \Rightarrow 0DDD \Rightarrow 01DD \Rightarrow 012D \Rightarrow 0127$  $N \Rightarrow ND \Rightarrow DD \Rightarrow 3D \Rightarrow 34$  $N \Rightarrow ND \Rightarrow NDD \Rightarrow DDD \Rightarrow 5DD \Rightarrow 56D \Rightarrow 568$ 

最右推导:

 $N \Rightarrow ND \Rightarrow N7 \Rightarrow ND7 \Rightarrow N27 \Rightarrow ND27 \Rightarrow N127 \Rightarrow D127 \Rightarrow 0127$  $N \Rightarrow ND \Rightarrow N4 \Rightarrow D4 \Rightarrow 34$  $N \Rightarrow ND \Rightarrow N8 \Rightarrow ND8 \Rightarrow N68 \Rightarrow D68 \Rightarrow 568$ 

## P36-7

 $G(S)$  $O \Rightarrow 1|3|5|7|9$  $N \Rightarrow 2|4|6|8|O$  $D \Rightarrow 0|N$  $S \Rightarrow O|AO$  $A \Rightarrow AD|N$ 

## P36-8

文法:

 $E \Rightarrow T|E \square T|E \square T$  $T \Rightarrow F|T * F|T / F$  $F \Rightarrow (E) i$ 

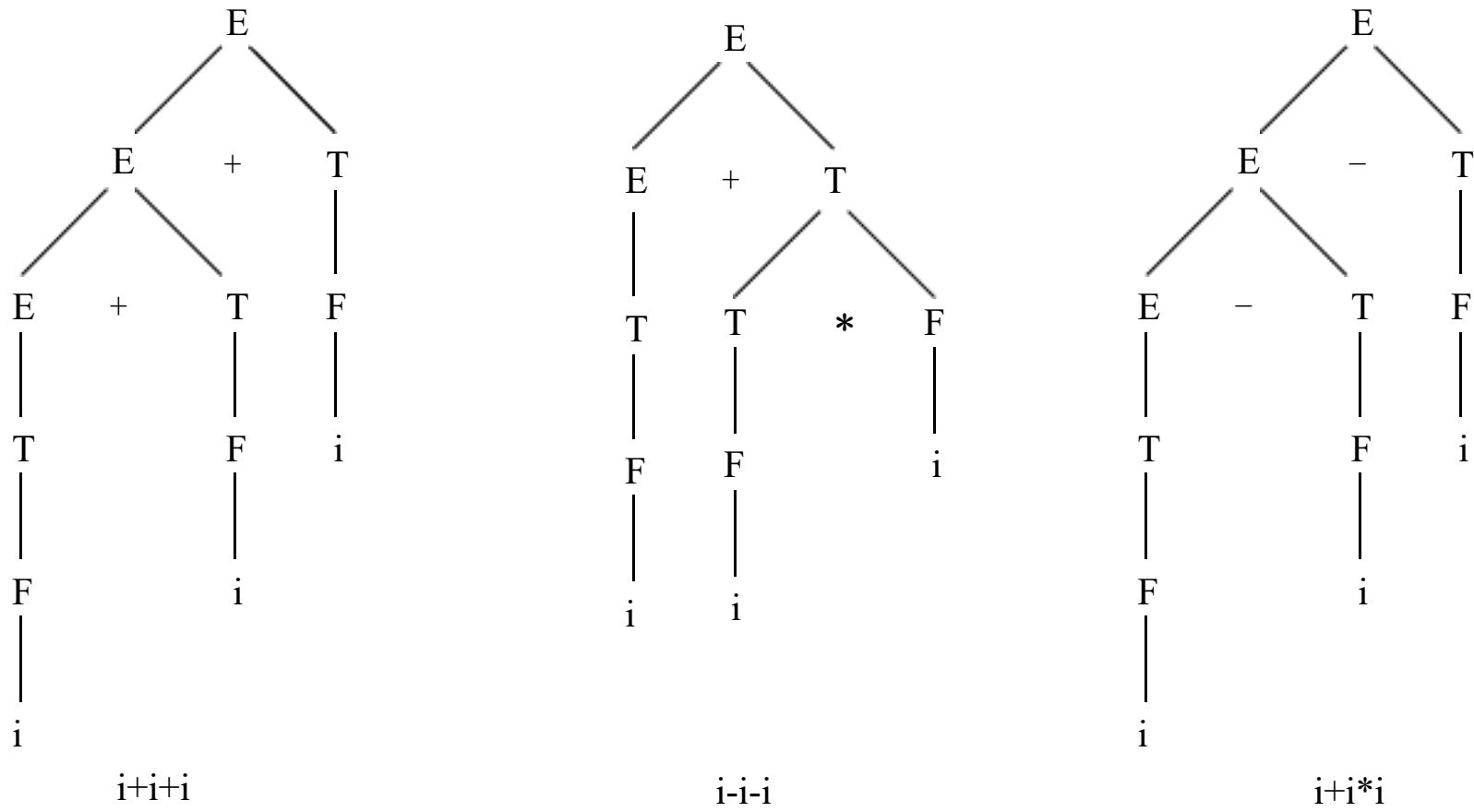
最左推导:

 $E \Rightarrow E \square T \square T \square T \square F \square T \square i \square T \square i \square T * F \square i \square F * F \square i \square i * F \square i \square i * i$  $E \Rightarrow T \square T * F \square F * F \square i * F \square i * (E) \square i * (E \square T) \square i * (T \square T) \square i * (F \square T)$  $\square i * (i \square T) \square i * (i \square F) \square i * (i \square i)$ 

最右推导:

$E \rightarrow E \mid T \mid E * F \mid E * i \mid E \mid F * i \mid E \mid i * i \mid T * i \mid F * i \mid i * i$   
 $E \rightarrow T \mid F * T \mid F * F \mid F * (E) \mid F * (E \mid T) \mid F * (E \mid F) \mid F * (E \mid )$   
 $\mid F * (T \mid ) \mid F * (F \mid ) \mid F * (i \mid ) \mid i * (i \mid )$

语法树: /\*\*\*\*\*



\*\*\*\*\*/

### P36-9

句子 iiei 有两个语法树:

$S \rightarrow iSeS \mid iSei \mid iiSei \mid iiei$   
 $S \rightarrow iS \mid iiSeS \mid iiSei \mid iiei$

### P36-10

/\*\*\*\*\*

$S \rightarrow TS \mid T$   
 $T \rightarrow (S) \mid ()$

\*\*\*\*\*/

### P36-11

/\*\*\*\*\*

L1:  
 $S \rightarrow AC$   
 $A \rightarrow aAb \mid ab$   
 $C \rightarrow cC \mid \epsilon$

L2:

S  $\square$  AB  
 A  $\square$  aA |  $\square$   
 B  $\square$  bBc | bc

L3:

S  $\square$  AB  
 A  $\square$  aAb |  $\square$   
 B  $\square$  aBb |  $\square$

L4:

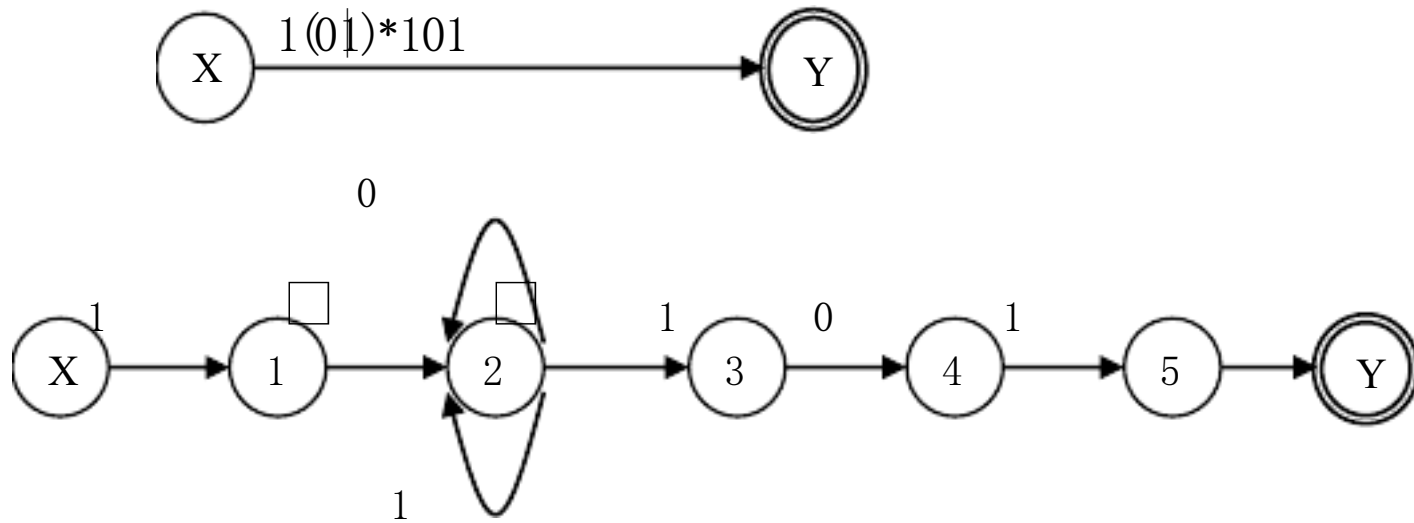
S  $\square$  A | B  
 A  $\square$  0A |  $\square$   
 B  $\square$  1B | A

\*\*\*\*\*/

### 第三章习题参考答案

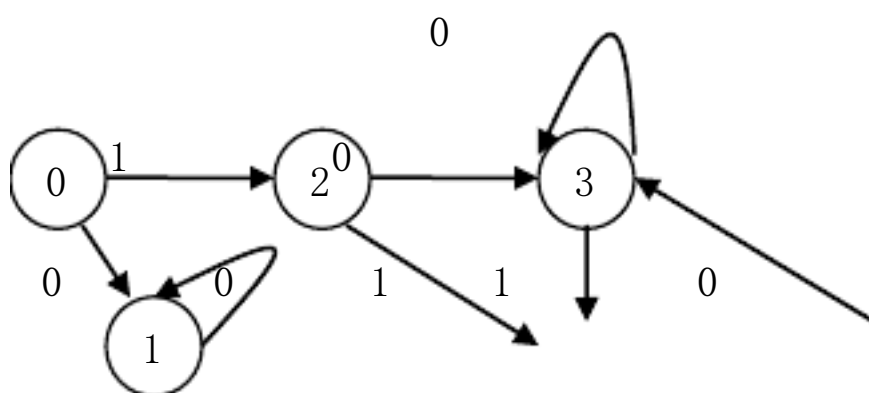
P64 -7

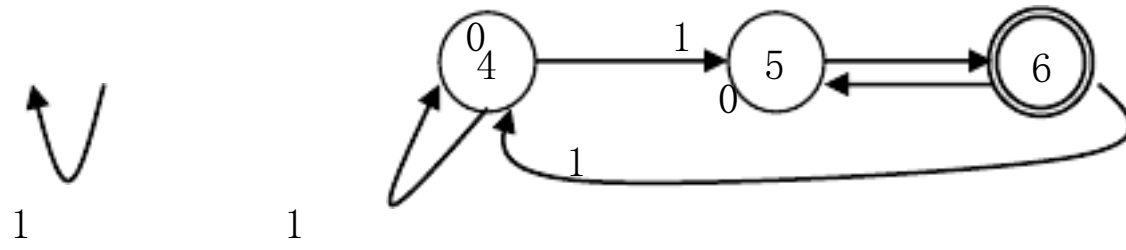
(1)



确定化:

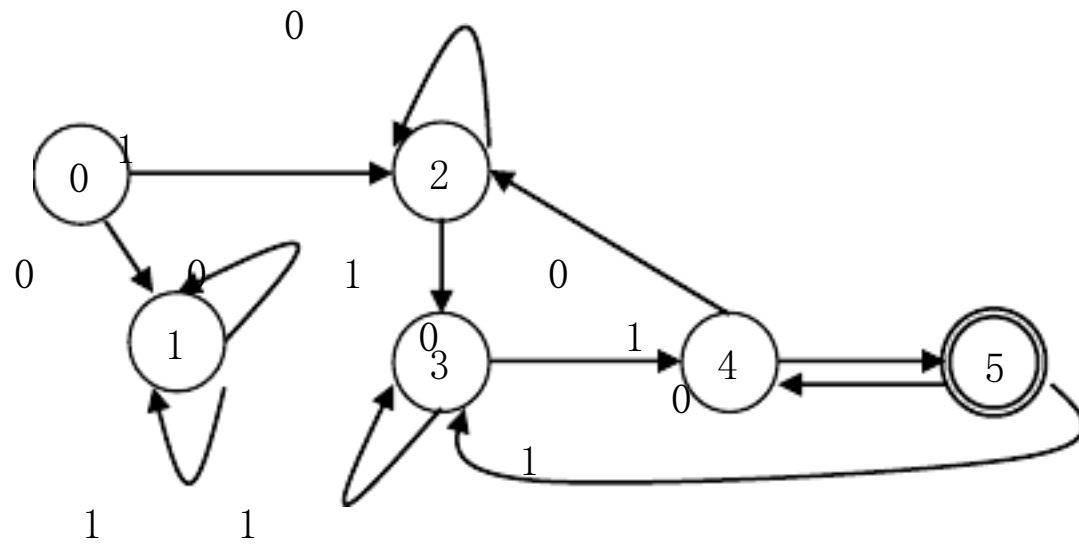
	0	1
{X}	$\phi$	{1, 2, 3}
$\phi$	$\phi$	$\phi$
{1, 2, 3}	{2, 3}	{2, 3, 4}
{2, 3}	{2, 3}	{2, 3, 4}
{2, 3, 4}	{2, 3, 5}	{2, 3, 4}
{2, 3, 5}	{2, 3}	{2, 3, 4, Y}
{2, 3, 4, Y}	{2, 3, 5}	{2, 3, 4, }





最小化:

- $\{0,1,2,3,4,5\}, 6$
- $\{0,1,2,3,4,5\} \square \{1,3,5\} \quad \{0,1,2,3,4,5\} \square \{1,2,4,6\}$
- $\{0,1,2,3,4\}, 5, 6$
- $\{0,1,2,3,4\} \square \{1,3,5\}$
- $\{0,1,2,3\}, 4, 5, 6$
- $\{0,1,2,3\} \square \{1,3\} \quad \{0,1,2,3\} \square \{1,2,4\}$
- $\{0,1\}, 2, 3, 4, 5, 6$
- $\{0,1\} \square \{1\} \quad \{0,1\} \square \{1,2\}$
- $\{2,3\}^0 \square \{3\} \quad \{2,3\}^1 \square \{4\}$
- $\{0\}, 1, 2, 3, 4, 5, 6$



P64 -8

(1)

$$(1 | 0)^* 01$$

(2)

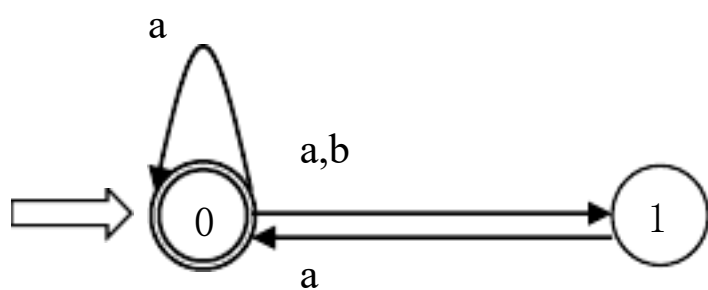
$$(1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9)^* (0 | 5) | (0 | 5)$$

(3)

$$0^* 1 (0 | 10^* 1)^* | 1^* 0 (0 | 10^* 1)^*$$

P64 -12

(a)

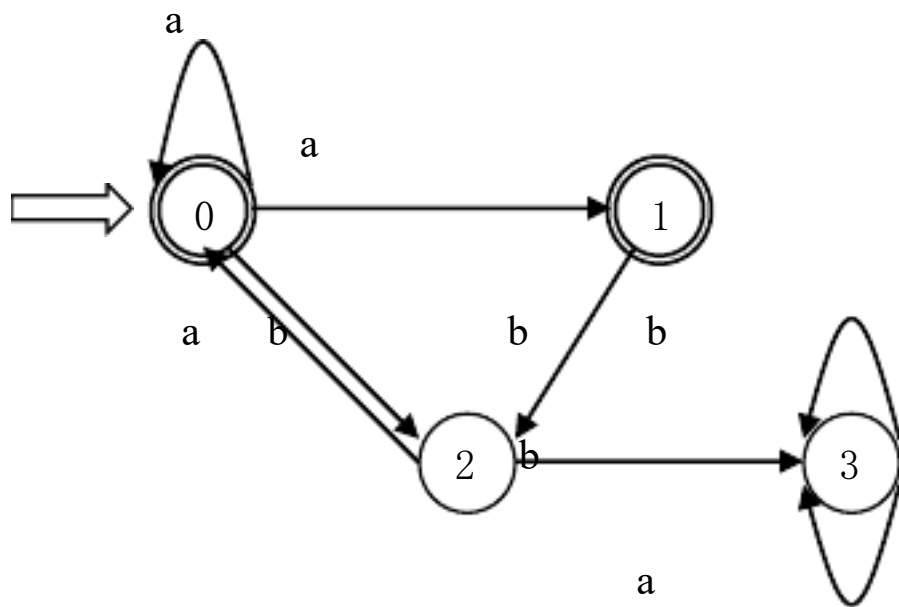


确定化:

	a	b
{0}	{0, 1}	{1}
{0, 1}	{0, 1}	{1}
{1}	{0}	$\phi$
$\phi$	$\phi$	$\phi$

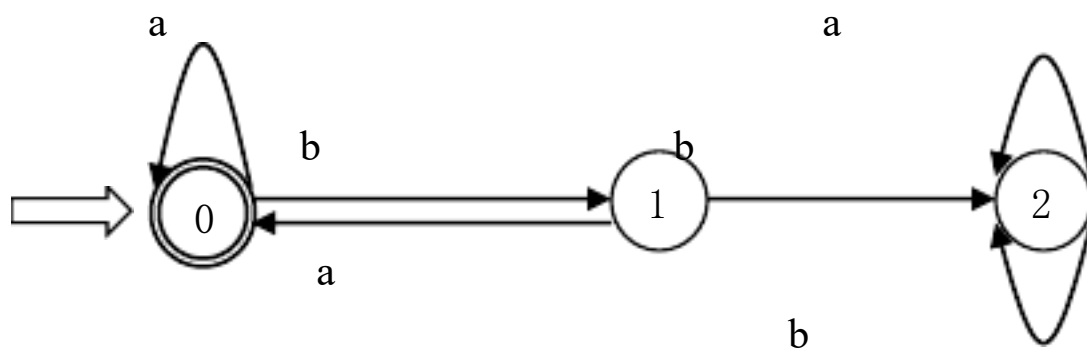
给状态编号:

	a	b
0	1	2
1	1	2
2	0	3
3	3	3

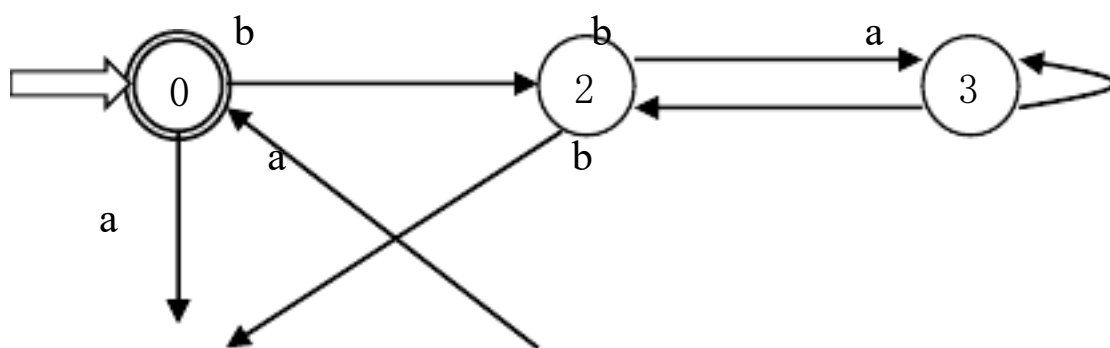


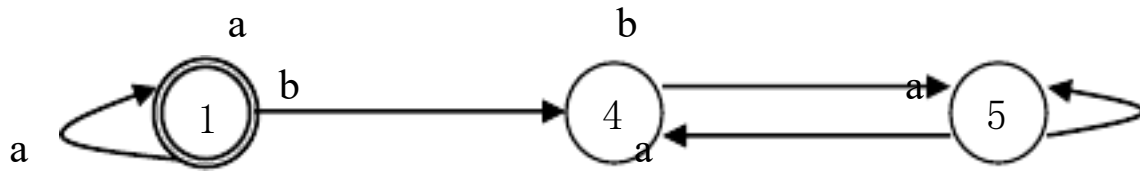
最小化:

$\{0,1\}, \{2,3\}$   
 $\{0,1\} \sqsubseteq \{1\} \quad \{0,1\} \sqsubseteq \{2\}$   
 $\{2,3\}^a \sqsubseteq \{0,3\} \quad \{2,3\}^b \sqsubseteq \{3\}$   
 $\{0,1\}^a, \{2\}, \{3\}$



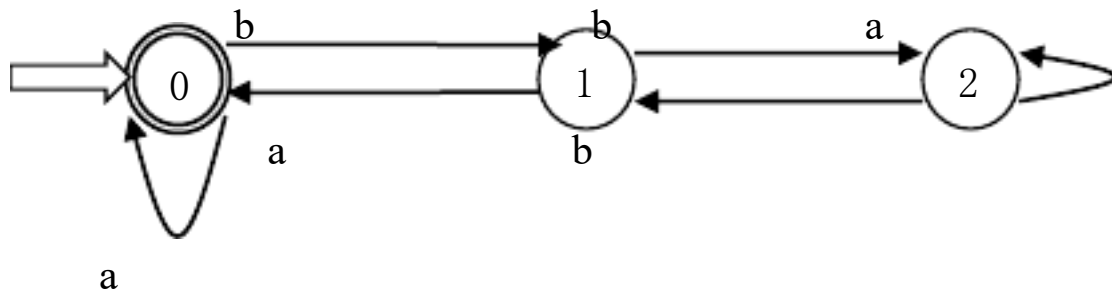
(b)



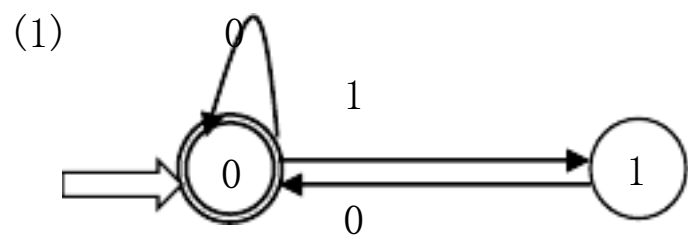


已经确定化了,进行最小化  
最小化:

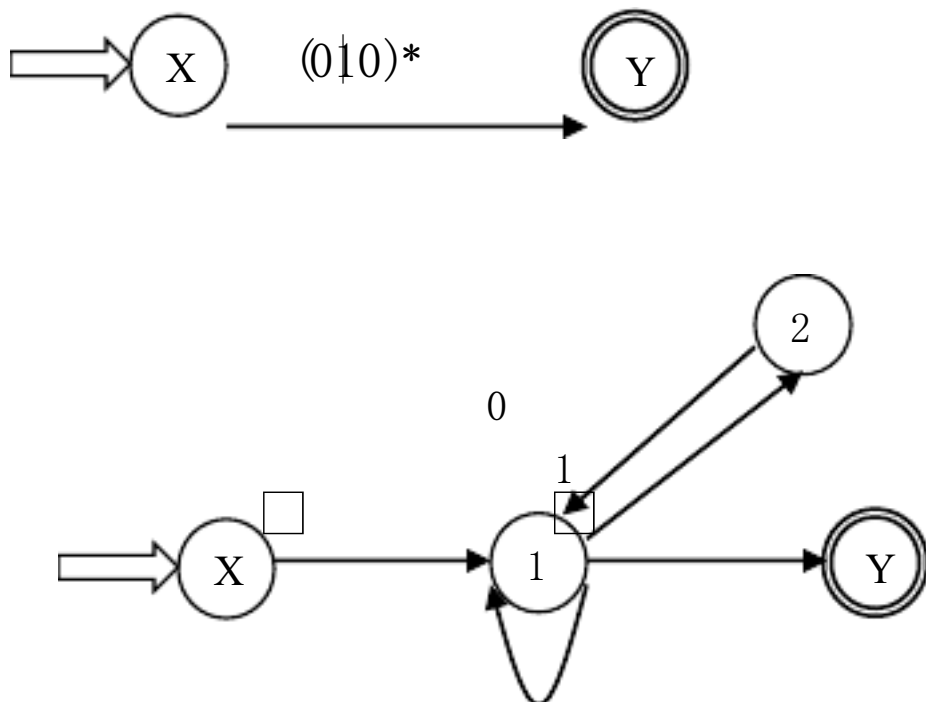
$\{\{0,1\}, \{2,3,4,5\}\}$   
 $\{0,1\} \stackrel{a}{\square} \{1\} \quad \{0,1\} \stackrel{b}{\square} \{2,4\}$   
 $\{2,3,4,5\} \stackrel{a}{\square} \{1,3,0,5\} \quad \{2,3,4,5\} \stackrel{b}{\square} \{2,3,4,5\}$   
 $\{2,4\} \stackrel{a}{\square} \{1,0\} \quad \{2,4\} \stackrel{b}{\square} \{3,5\}$   
 $\{3,5\} \stackrel{a}{\square} \{3,5\} \quad \{3,5\} \stackrel{b}{\square} \{2,4\}$   
 $\{\{0,1\}, \{2,4\}, \{3,5\}\}$   
 $\{0,1\} \stackrel{a}{\square} \{1\} \quad \{0,1\} \stackrel{b}{\square} \{2,4\}$   
 $\{2,4\} \stackrel{a}{\square} \{1,0\} \quad \{2,4\} \stackrel{b}{\square} \{3,5\}$   
 $\{3,5\} \stackrel{a}{\square} \{3,5\} \quad \{3,5\} \stackrel{b}{\square} \{2,4\}$



P64 -14



(2):



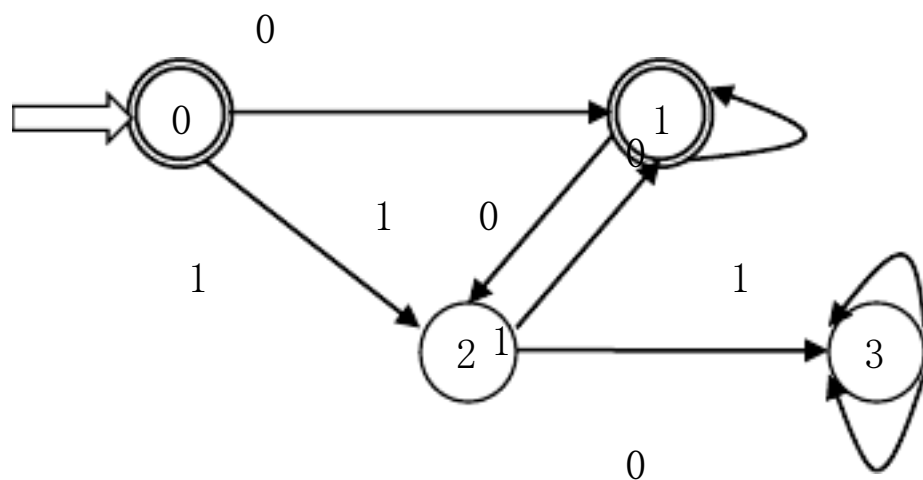
0

确定化:

	0	1
{X,1,Y}	{1,Y}	{2}
{1,Y}	{1,Y}	{2}
{2}	{1,Y}	$\phi$
$\phi$	$\phi$	$\phi$

给状态编号:

	0	1
0	1	2
1	1	2
2	1	3
3	3	3



最小化:

$\{0,1\}, \{2,3\}$

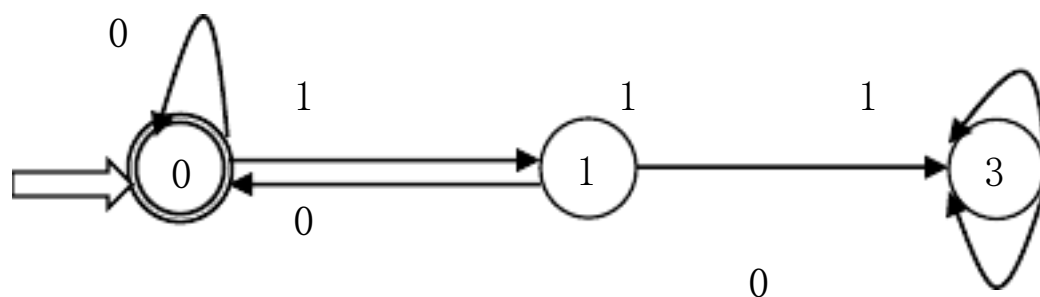
$\{0,1\} \stackrel{0}{\sqsubset} \{1\}$

$\{0,1\} \stackrel{1}{\sqsubset} \{2\}$

$\{2,3\} \stackrel{0}{\sqsubset} \{1,3\}$

$\{2,3\} \stackrel{1}{\sqsubset} \{3\}$

$\{0,1\}, \{2\}, \{3\}$



### 第四章

P81 -1

(1) 按照 T,S 的顺序消除左递归

$G[S]$

$S \rightarrow a \mid \wedge(T)$

$T \rightarrow ST \mid \epsilon$

$T \rightarrow \square, ST \rightarrow \square \square$

递归子程序:

```

procedure S;
begin
  if sym='a' or sym='^'
  then advance
  else if sym='('
  then begin
    advance;T;
    if sym=')' then advance;
    else error;
  end
  else error
end;
procedure T;
begin
  S;T□
end;
procedure T□;
begin
  if sym=','
  then begin
    advance;
    S;T□
  end
end;

```

其中:

sym:是输入串指针 IP 所指的符号

advance:是把 IP 调至下一个输入符号

error:是出错诊察程序

(2)

$FIRST(S) = \{a, ^, (\}$

$FIRST(T) = \{a, ^, (\}$

$FIRST(T□) = \{, □\}$

$FOLLOW(S) = \{), ,, \#\}$

$FOLLOW(T) = \{\}$

$FOLLOW(T□) = \{\}$

预测分析表

	a	^	(	)	,	#
S	S □ a	S □ ^	S □ (T)			
T	T □ ST □	T □ ST □	T □ ST □			
T□				T □ □ □	T □ □ □ ,ST □	

是 LL(1)文法

P81 -2

文法:



$E \rightarrow TE$   
 $E \rightarrow \epsilon$   
 $T \rightarrow FT$   
 $T \rightarrow \epsilon$   
 $F \rightarrow PF$   
 $F \rightarrow *F$   
 $P \rightarrow (E) | a | b | \wedge$

(1)

$FIRST(E) = \{(a,b,\wedge)\}$   
 $FIRST(E') = \{+, \epsilon\}$   
 $FIRST(T) = \{(a,b,\wedge)\}$   
 $FIRST(T') = \{(a,b,\wedge), \epsilon\}$   
 $FIRST(F) = \{(a,b,\wedge)\}$   
 $FIRST(F') = \{*, \epsilon\}$   
 $FIRST(P) = \{(a,b,\wedge)\}$   
 $FOLLOW(E) = \{#, \wedge\}$   
 $FOLLOW(E') = \{#, \wedge\}$   
 $FOLLOW(T) = \{+, \wedge, \#\}$   
 $FOLLOW(T') = \{+, \wedge, \#\}$   
 $FOLLOW(F) = \{(a,b,\wedge, +, \wedge), \#\}$   
 $FOLLOW(F') = \{(a,b,\wedge, +, \wedge), \#\}$   
 $FOLLOW(P) = \{*, (a,b,\wedge, +, \wedge), \#\}$

(2)

考虑下列产生式:

$E \rightarrow E$   
 $T \rightarrow T$   
 $F \rightarrow *F$   
 $P \rightarrow (E) | \wedge | a | b$

$FIRST(+E) \cap FIRST(\epsilon) = \{+\} \cap \{\epsilon\} = \phi$   
 $FIRST(+E) \cap FOLLOW(E') = \{+\} \cap \{#, \wedge\} = \phi$   
 $FIRST(T) \cap FIRST(\epsilon) = \{(a,b,\wedge)\} \cap \{\epsilon\} = \phi$   
 $FIRST(T) \cap FOLLOW(T') = \{(a,b,\wedge)\} \cap \{+, \wedge, \#\} = \phi$   
 $FIRST(*F) \cap FIRST(\epsilon) = \{*\} \cap \{\epsilon\} = \phi$   
 $FIRST(*F) \cap FOLLOW(F') = \{*\} \cap \{(a,b,\wedge, +, \wedge), \#\} = \phi$   
 $FIRST((E)) \cap FIRST(a) \cap FIRST(b) \cap FIRST(\wedge) = \phi$

所以,该文法为 LL(1)文法.

(3)

	+	*	(	)	a	b	$\wedge$	#
E			$E \rightarrow TE$		$E \rightarrow TE$	$E \rightarrow TE$	$E \rightarrow TE$	
E'	$E \rightarrow E$			$E \rightarrow \epsilon$				$E \rightarrow \epsilon$
T			$T \rightarrow FT$		$T \rightarrow FT$	$T \rightarrow FT$	$T \rightarrow FT$	

--	--	--	--	--	--	--	--	--

	$\square \square$		$T \square \square T$	$T \square \square \square$	$T \square \square T$	$T \square \square T$	$T \square \square T$	$T \square \square \square$
F			$F \square PF \square$		$F \square PF \square$	$F \square PF \square$	$F \square PF \square$	
F'	$F \square \square \square$	$F \square \square * F \square$	$F \square \square \square$	$F \square \square \square$	$F \square \square \square$	$F \square \square \square$	$F \square \square \square$	$F \square \square \square$
P			$P \square (E)$		$P \square a$	$P \square b$	$P \square ^$	

(4)

procedure E;

begin

if sym='(' or sym='a' or sym='b' or sym='^'

then begin T; E' end

else error

end

procedure E';

begin

if sym='+'

then begin advance; E end

else if sym <> ')' and sym <> '#' then error

end

procedure T;

begin

if sym='(' or sym='a' or sym='b' or sym='^'

then begin F; T' end

else error

end

procedure T';

begin

if sym='(' or sym='a' or sym='b' or sym='^'

then T

else if sym='\*' then error

end

procedure F;

begin

if sym='(' or sym='a' or sym='b' or sym='^'

then begin P; F' end

else error

end

procedure F';

begin

if sym='\*'

then begin advance; F' end

end

procedure P;

```

if sym='a' or sym='b' or sym='^'
  then advance
  else if sym='(' then
    begin
      advance; E;
      if sym=')' then advance
      else error
    end
  else error
end;

```

3

/\*\*\*\*\*

- (1)
- (2) 不是，对于 A 不满足条件 3。
- (3) 不是，A、B 均不满足条件 3。
- (4) 是，满足三个条件。

\*\*\*\*\*/

## 第五章

P133 -1

$E \rightarrow T \mid E \rightarrow T * F$

短语:  $E+T * F, T * F,$

直接短语:  $T * F$

句柄:  $T * F$

P133 -2

文法:

$S \rightarrow a \mid \hat{\ } \mid (T)$   
 $T \rightarrow T, S \mid S$

(1)

最左推导:

$S \rightarrow (T) \rightarrow (T, S) \rightarrow (S, S) \rightarrow (a, S) \rightarrow (a, (T)) \rightarrow (a, (T, S)) \rightarrow (a, (S, S)) \rightarrow (a, (a, S)) \rightarrow (a, (a, a))$   
 $S \rightarrow (T, S) \rightarrow (S, S) \rightarrow ((T), S) \rightarrow ((T, S), S) \rightarrow ((T, S, S), S) \rightarrow ((S, S, S), S) \rightarrow (((T), S, S), S)$   
 $\rightarrow (((T, S), S, S), S) \rightarrow (((S, S), S, S), S) \rightarrow (((a, S), S, S), S) \rightarrow (((a, a), S, S), S)$   
 $\rightarrow (((a, a), \hat{\ }, S), S) \rightarrow (((a, a), \hat{\ }, (T)), S) \rightarrow (((a, a), \hat{\ }, (S)), S) \rightarrow (((a, a), \hat{\ }, a)) S$   
 $\rightarrow (((a, a), \hat{\ }, a)) a$

最右推导:

$T) \square (T,S) \square (T, \underline{T}) \square (T, \underline{T},S) \square (T, \underline{T},a) \square (T, \underline{S},a) \square (T, \underline{a},a)$   
 $\square (S, \underline{a},a) \square (a, \underline{a},a)$   
 $S \square (T,S) \square (T,a) \square (S,a) \square (\underline{T},a) \square (\underline{T},S,a) \square (\underline{T}, \underline{T})a \square (\underline{T}, \underline{S})a$   
 $\square (\underline{T}, \underline{a})a \square (\underline{T},S, \underline{a})a \square (\underline{T}, \hat{\quad}, \underline{a})a \square (\underline{S}, \hat{\quad}, \underline{a})a \square ((\underline{T}), \hat{\quad}, \underline{a})a$   
 $\square ((\underline{T},S), \hat{\quad}, \underline{a})a \square ((\underline{T},a), \hat{\quad}, \underline{a})a \square ((\underline{S},a), \hat{\quad}, \underline{a})a \square ((\underline{a},a), \hat{\quad}, \underline{a})a$

(2)

$((\underline{a}, \hat{\quad}, (a)), a)$   
 $((\underline{S}, a), \hat{\quad}, (a)), a)$   
 $((\underline{T}, a), \hat{\quad}, (a)), a)$   
 $((\underline{T}, S), \hat{\quad}, (a)), a)$   
 $((\underline{T}), \hat{\quad}, (a)), a)$   
 $((\underline{S}), \hat{\quad}, (a)), a)$   
 $((T, \hat{\quad}, (a)), a)$   
 $((\underline{T}, S), (a)), a)$   
 $((T, (\underline{a})), a)$   
 $((T, (\underline{S})), a)$   
 $((T, (\underline{T})), a)$   
 $((\underline{T}, S), a)$   
 $((\underline{T}), a)$   
 $((\underline{S}), a)$   
 $((T, S))$   
 $((T))$   
 $S$

-归约”过程:

步骤	栈	输入串	动作
0	#	$((\underline{a}, a), \hat{\quad}, (a)), a) \#$	预备
1	#(	$(\underline{a}, a), \hat{\quad}, (a)), a) \#$	进
2	#((	$(\underline{a}, a), \hat{\quad}, (a)), a) \#$	进
3	#(((	$\underline{a}, a), \hat{\quad}, (a)), a) \#$	进
4	#(((a	$, a), \hat{\quad}, (a)), a) \#$	进
5	#(((S	$, a), \hat{\quad}, (a)), a) \#$	归
6	#(((T	$, a), \hat{\quad}, (a)), a) \#$	归
7	#(((T,	$a), \hat{\quad}, (a)), a) \#$	进
8	#(((T,a	$), \hat{\quad}, (a)), a) \#$	进
9	#(((T,S	$), \hat{\quad}, (a)), a) \#$	归
10	#(((T	$), \hat{\quad}, (a)), a) \#$	归
11	#(((T	$, \hat{\quad}, (a)), a) \#$	进
12	#((S	$, \hat{\quad}, (a)), a) \#$	归
13	#((T	$, \hat{\quad}, (a)), a) \#$	归
14	#((T,	$\hat{\quad}, (a)), a) \#$	进
15	#((T,^	$, (a)), a) \#$	进
16	#((T,S	$, (a)), a) \#$	归
17	#((T	$, (a)), a) \#$	归

$\#((T, (a)),a)\#$   
 19  $\#((T,( a)),a)\#$  进  
 20  $\#((T,(a )),a)\#$  进  
 21  $\#((T,(S )),a)\#$  归  
 22  $\#((T,(T )),a)\#$  归  
 23  $\#((T,(T )),a)\#$  进  
 24  $\#((T,S ),a)\#$  归  
 25  $\#((T ),a)\#$  归  
 26  $\#((T ),a)\#$  进  
 27  $\#(S ,a)\#$  归  
 28  $\#(T ,a)\#$  归  
 29  $\#(T ,a)\#$  进  
 30  $\#(T,a )\#$  进  
 31  $\#(T,S )\#$  归  
 32  $\#(T )\#$  归  
 33  $\#(T )\#$  进  
 34  $\#S \#$  归

3

(1)

$FIRSTVT(S) = \{a, ^, (\}$

$FIRSTVT(T) = \{., a, ^, (\}$

$LASTVT(S) = \{a, ^, )\}$

$LASTVT(T) = \{., a, ^, )\}$

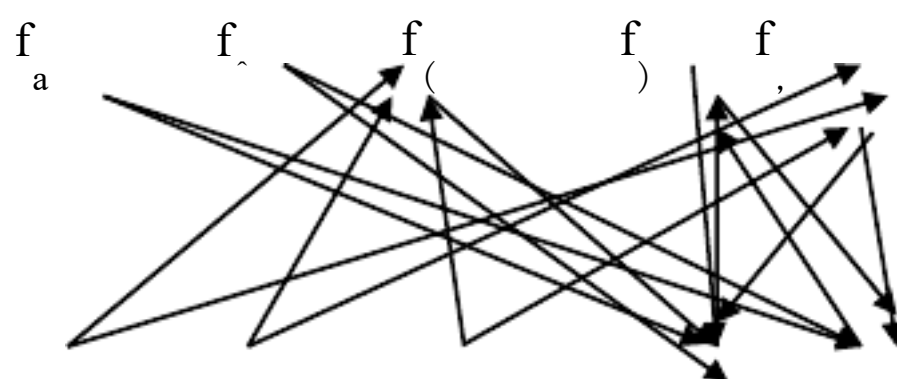
(2)

	a	^	(	)	,
a				>	>
^				>	>
(	<	<	<	=	<
)				>	>
,	<	<	<	>	>

是算符文法，并且是算符优先文法

(3) 优先函数

	a	^	(	)	,
f	4	4	2	4	4
g	5	5	5	2	3



a g ( g ) g ,

4)

栈	输入字符串	动作
#	(a,(a,a)) #	预备
#(	a, (a,a)#	进
#(a	, (a,a)#	进
#(t	, (a,a)#	归
# (t,	(a,a) #	进
# (t, (	a,a) #	进
# (t, (a	,a) #	进
# (t, (t	,a) #	归
# (t, (t,	a) #	进
# (t, (t,a	) #	进
# (t, (t,s	) #	归
# (t, (t	) #	归
# (t, (t)	) #	进
# (t,s	) #	归
# (t	) #	归
# (t)	#	进
#s	#	归

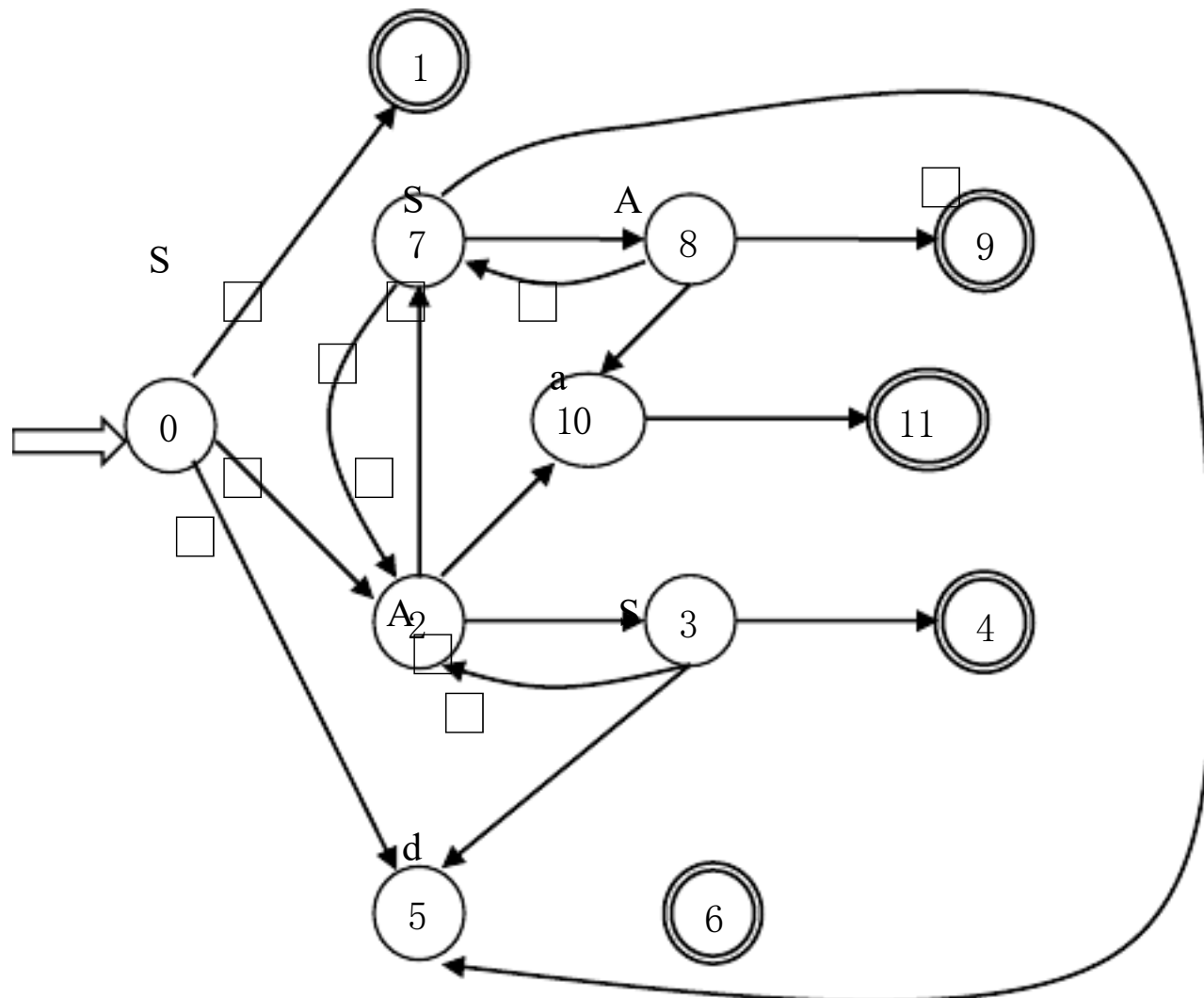
success

5

(1)

0. S □ S      1. S □ □ S □      2. S □ AS      3. S □ AS
4. S □ AS □      5. S □ b      6. s □ b □      7. A □ SA
8. A □ S □ A      9. A □ SA □      10. A □ a      11. A □ a □

(2)



以上内容仅为本文档的试下载部分，为可阅读页数的一半内容。如要下载或阅读全文，请访问：<https://d.book118.com/228134033005006024>